

# Electrical Engineering 229A Lecture 26 Notes

Daniel Raban

November 30, 2021

## 1 Convex Dual of the Cumulant Generating Function and Sanov's Theorem

### 1.1 The cumulant generating function and convex duality

Suppose  $X \in \mathbb{R}^d$  is a random variable.

**Definition 1.1.** The map  $\theta \mapsto \mathbb{E}[e^{\theta^\top X}]$  with  $\theta \in \mathbb{R}^d$  is called the **moment generating function**.

**Definition 1.2.** The map  $\theta \mapsto \log \mathbb{E}[e^{\theta^\top X}]$  with  $\theta \in \mathbb{R}^d$  is called the **cumulant generating function**.

If we differentiate the moment generating function with respect to  $\theta$  and set  $\theta = 0$ , we get the moments of  $X$ . Likewise, doing the same to the cumulant generating function gives us the cumulants of  $X$ . One advantage of working with the cumulant generating function is that it is convex.

We have dealt with finite (and countable) random variables and some densities. For a finite random variable  $X \in \mathcal{X}$  with  $|\mathcal{X}| = d$ , it is interesting to consider  $Z \in \mathbb{R}^d$  where  $Z = e_i$  with probability  $p_i$  (here,  $e_i$  is the  $i$ -th basis vector). Then

$$\log \mathbb{E}[e^{\theta^\top Z}] = \log \sum_{i=1}^d p_i e^{\theta_i}$$

because  $\theta^\top e_i = \theta_i$  for  $i = 1, \dots, d$ .

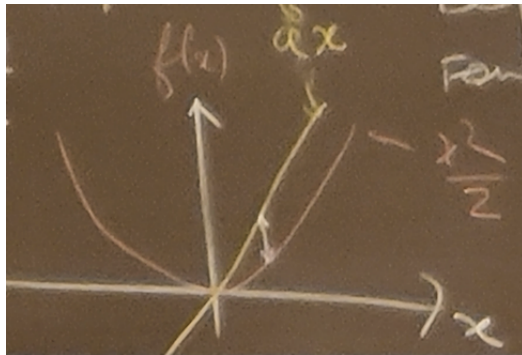
To any (extended real-valued) convex function there is a *dual*<sup>1</sup> convex function on  $\mathbb{R}^d$ .

**Example 1.1.** Let  $d = 1$  and consider  $f(x) = x^2/2$ . Consider a line  $ax$  of slope  $a$  and look at the height that separates the line from the function. Find the point at which this

---

<sup>1</sup>This is sometimes called Fenchel duality, Legendre duality, or Fenchel-Legendre duality.

height is the greatest to calculate the dual  $\hat{f}(a) := \sup_{x \in \mathbb{R}} ax - f(x)$ .



Here, we can calculate  $\hat{f}(a) = a^2/2$ . In a related sense to how the Gaussian is self-dual for the Fourier transform, this function is self-dual for the Fenchel-Legendre transform.

**Example 1.2.** Let  $f(x) = e^x$ . To find  $\hat{f}(a)$ , since  $f'(x) = e^x = a$  for  $x$ , if  $a > 0$ , this occurs if  $x = \ln a$ , and if  $a \leq 0$ , this is impossible. So we get

$$\begin{aligned} \hat{f}(a) &= \sup_x (ax - e^x) \\ &= \begin{cases} a \ln a - a & a > 0 \\ 0 & a = 0 \\ \infty & a < 0. \end{cases} \end{aligned}$$

What if  $d > 1$ ?

**Definition 1.3.** Suppose  $\Phi : \mathbb{R}^d \mapsto \mathbb{R} \cup \{\infty\}$  is convex. Its **Fenchel-Legendre dual** is

$$\hat{\Phi}(a) := \sup_{x \in \mathbb{R}^d} a^\top x - \Phi(x)$$

for  $a \in \mathbb{R}^d$ .

Again,

$$\hat{\Phi}(a) = a^\top x_a - \Phi(x_a),$$

where  $x_a$  is defined by  $\nabla \Phi(x_a) = a$  (if  $x_a$  exists). It can be shown that

$$\Phi(x) = \sup_a x^\top a - \hat{\Phi}(a).$$

To check this where  $\Phi$  expresses all derivatives, write

$$\Phi(x) \geq x^\top a - \hat{\Phi}(a) \quad \forall x, a \iff \hat{\Phi}(a) \geq a^\top x - \Phi(x) \quad \forall x, a.$$

**Proposition 1.1.** Let  $X$  take values in  $\mathcal{X}$  with  $|\mathcal{X}| = d$  and  $p_i = \mathbb{P}(X = i)$ . Let  $Z = e_i$  iff  $X = i$  (i.e.  $P(Z = e_i) = p_i$  for  $1 \leq i \leq d$ ). Then the Fenchel dual of  $\Phi(\theta) = \ln \mathbb{E}[e^{\theta^\top Z}]$  is

$$\widehat{\Phi}(a) = \begin{cases} D(a \parallel p) & \text{if } a \text{ is a probability distribution} \\ \infty & \text{otherwise.} \end{cases}$$

*Proof.* Here,

$$\Phi_Z(\theta) = \ln \sum_{i=1}^d p_i e^{\theta_i},$$

so

$$\nabla \Phi_Z(\theta) = \begin{bmatrix} \frac{p_1 e^{\theta_1}}{\sum_{i=1}^d p_i e^{\theta_i}} \\ \vdots \end{bmatrix}.$$

This expresses only gradients that are probability distributions (means where  $p_i \neq 0$ ). We have

$$\widehat{\Phi}_X(a) = a^\top p_a - \ln \sum_{i=1}^d p_i e^{\theta_{a_i}},$$

where  $\theta_a$  is defined in terms of  $a$  via  $\nabla \Phi(\theta_a) = a$ , i.e.  $p_i e^{\theta_i}$  is proportional to  $a_i$  (i.e.  $\theta_i = \ln \frac{a_i}{p_i} + \text{constant}$ ). The constant is  $\log \sum_{i=1}^d p_i e^{(\theta_a)_i} = 0$ .

$$\begin{aligned} &= \sum_{i=1}^d a_i \ln \frac{a_i}{p_i} - \ln \left( \sum_{i=1}^d p_i e^{\ln \frac{a_i}{p_i}} \right) \\ &= D(a \parallel p). \end{aligned} \quad \square$$

## 1.2 Large deviations and Sanov's theorem

Roughly speaking, a basic large deviations theory result is of the form: If  $Z_1, Z_2, \dots$  are iid  $\mathbb{R}^d$ -valued with  $\log \mathbb{E}[e^{\theta^\top Z}]$  denoted  $\Phi_Z(\theta)$  and  $\mathbb{E}[Z_1] = 0 \in \mathbb{R}^d$ , then for any open set  $A \subseteq \mathbb{R}^d$ ,

$$\liminf_{n \rightarrow \infty} -\frac{1}{n} \log \mathbb{P} \left( \frac{Z_1 + \dots + Z_n}{n} \in A \right) \leq \inf_{z \in A} \widehat{\Phi}_Z(z).$$

Here is a special case.

If  $X_1, X_2, \dots$ , are i.i.d.  $\mathcal{X}$ -valued with  $\mathcal{X} = \{1, 2, \dots, d\}$  and  $Z_1, Z_2, \dots$  are i.i.d.  $\mathbb{R}^d$ -valued created from  $X_1, X_2, \dots$ , then observe that  $\frac{Z_1 + \dots + Z_n}{n}$  is equivalent to the empirical distribution of  $(X_1, \dots, X_n)$ , i.e.  $\frac{Z_1 + \dots + Z_n}{n} = \sum_{i=1}^d \frac{N(i|x^n)}{n} e_i$ . Let  $P_{x^n} := (\frac{N(i|x^n)}{n}, i = 1, \dots, d)$ . So for any open subset  $A \subseteq \text{simplex in } \mathbb{R}^d$ ,

$$\liminf_n -\frac{1}{n} \log \mathbb{P}(P_{X^n} \in A) \leq \inf_{a \in A} D(a \parallel p).$$

Recall that if  $x^n = (x_1, \dots, x_n) \in \mathcal{X}^n$  with  $\mathcal{X}$  finite and if  $\mathcal{P}$  denotes the set of probability distributions on  $X$ , then  $p_{x^n} \in \mathcal{P}$  denotes  $(\frac{N(x|x^n)}{n}, x \in \mathcal{X})$  and  $\mathcal{P}_n$  denotes the set of all such  $P_{x^n}$ . For an  $n$ -**type**  $P \in \mathcal{P}_n$ , the **typicality set for  $P$**  refers to  $T(P) := \{x^n \in \mathcal{X}^n : P_{x^n} = P\}$ . For  $Q \in \mathcal{P}$ ,

$$\begin{aligned} Q(x^n) &= \prod_{i=1}^n q(x_i) \\ &= \prod_{x \in X} q(x)^{N(x|x^n)} \\ &= 2^{-n(H(P_{x^n}) + D(P_{x^n} \| Q))}. \end{aligned}$$

We also proved that for  $P \in \mathcal{P}_n$ ,

$$P^n(T(P)) \geq P^n(T(\tilde{P})) \quad \forall \tilde{P} \in \mathcal{P}_n,$$

$|\mathcal{P}_n| \leq (n+1)^{|\mathcal{X}|}$ , and for  $P \in \mathcal{P}_n$ ,

$$\frac{1}{(n+1)^{|\mathcal{X}|}} 2^{nH(P)} \leq |T(P)| \leq 2^{nH(P)}.$$

**Theorem 1.1** (Sanov). *Let  $\mathcal{X}$  be finite,  $X_1, X_2, \dots \stackrel{\text{iid}}{\sim} Q$ , and  $E \subseteq \mathcal{P}$ . Assume that  $E$  is the closure of its interior. Then*

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log Q^n(P_{X^n} \in E) = -D(P^* \| Q),$$

where

$$P^* = \arg \min_{P \in E} D(P \| Q).$$

**Remark 1.1.** Since  $E$  is closed and  $D(\cdot \| Q)$  is continuous, this argmin exists.  $P^*$  is called the *I-projection* of  $Q$  onto  $E$ .

*Proof.* For the upper bound,

$$\begin{aligned} Q^n(P_{X^n} \in E) &= Q^n(P_{X^n} \in E \cap \mathcal{P}_n) \\ &\leq (n+1)^{|\mathcal{X}|} 2^{-nD(P^* \| Q)} \end{aligned}$$

For the lower bound, for any  $P^{(n)} \in \mathcal{P}_n \cap E$ ,

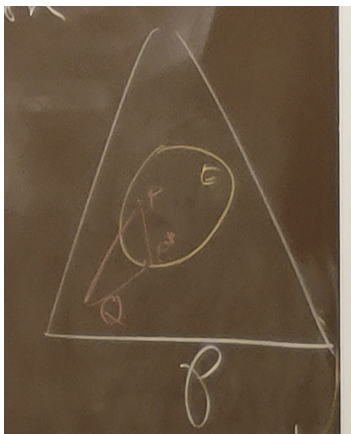
$$\begin{aligned} Q^n(P_{X^n} \in E) &\geq Q^n(T(P^{(n)})) \\ &\geq \frac{1}{(n+1)^{|\mathcal{X}|}} 2^{-nD(P^{(n)} \| Q)}. \end{aligned}$$

Choose  $P^{(n)} \rightarrow P^*$ . □

Here is a nice observation about the  $I$ -projection of  $Q$  onto a *convex* set  $E$ .

**Proposition 1.2.** For all  $P \in E$ ,

$$D(P \parallel Q) \geq D(P \parallel P^*) + D(P^* \parallel Q).$$



This tells us that we should think of  $D(P \parallel Q)$  as the *square* of a distance.

*Proof.* Consider the relative entropy  $D(\lambda P + (1 - \lambda)P^* \parallel Q)$  for  $\lambda \in [0, 1]$ . Differentiate in  $\lambda$ . It must be nonnegative.  $\square$